

No calculators or books. Submit your answers on a separate paper.

1. Consider the following algorithm. Assume  $A[1\dots n]$  is an array that contains integers.

```

1  STUPID( $A, p, q$ )
2  if  $p \geq q$  return
3   $i := p + 1$ 
4  while  $i < q$  do
5      if  $A[i] < A[p]$  then
6          SWAP( $A[i], A[p]$ )
7      endif
8      if  $A[i] > A[q]$  then
9          SWAP( $A[i], A[q]$ )
10     endif
11      $i := i + 1$ 
12 endwhile
    STUPID( $A, p + 1, q - 1$ )

```

- (a) (2 points) Write a recurrence equation for the time consumption of DUMMY( $A, 1, n$ ).
- (b) (2 points) What does the algorithm do to the array  $A$ ?
- (c) (4 points) Change the algorithm so that it is no longer recursive. Give invariants to both loops in the iterative version.
2. Give the  $\Theta$ , or  $O$  and  $\Omega$ - class of the following functions. Be as accurate as possible.

(a) (2 points)  $\sum_{i=1}^n \sqrt{i}$

(b) (2 points)  $\log(n!)$

(c) (2 points)  $\sum_{i=1}^{\lfloor \log n \rfloor} i$

(d) (2 points)  $O(n^2) + \Theta(n \log n)$

(e) (2 points)  $n^2 \log n + n^3$

3. Solve the following recurrences, assumin  $T(n)$  is a constant for  $n < 2$ . Simply give the answer in  $\Theta$ .

(a) (2 points)  $T(n) = T(n - 1) + n/2$

(b) (2 points)  $T(n) = T(n/2) + n$

(c) (2 points)  $T(n) = 3T(n/3) + n$

(d) (2 points)  $T(n) = 4T(n/2) + n$

(e) (2 points)  $T(n) = T(\sqrt{n}) + \sqrt{n}$

4. Let  $(V, E)$  be a directed graph. We wish to represent the graph by either using adjacency matrix or adjacency list. A list requires 32 bits for each vertex and 64 bits for each edge. A matrix, on the other hand, requires only 8 bits for each slot.

- (a) (4 points) Assuming you wish to minimize the amount of memory needed for your graph, how does your choice (matrix or list) depend on the number of edges in the graph?
- (b) (4 points) We are solving a problem about graph, and the problem can be solved by an algorithm, that looks like the following. ( $Q$  is a priority queue.)

```

1   F(V, E)
2   Q := V
3   ...
4   while Q ≠ ∅ do
5     v := Q.EXTRACT-MIN()
6     for u ∈ V
7       if (u, v) ∈ E
8         ...
9     endif
10    endfor
11  endwhile

```

Assume that everything that is omitted (marked by  $\dots$ ) is constant time, and the assignment  $Q := V$  is  $\Theta(V)$ , and  $Q.EXTRACT-MIN()$  is  $\Theta(\log n)$  when  $Q$  contains  $n$  elements, and nothing is ever added to  $Q$ . Analyse the time consumption for this algorithm for both matrices and lists.

5. In the following assignment,  $T_1$  and  $T_2$  are red-black trees and  $A$  is an array. Assume emptying a tree is constant time.  $T.(max)()$  returns the largest element in the tree.

```

1   DUMMY2(A[1, ..., n])
2   if n ≤ 1 then return
3   T1 := ∅
4   T2 := ∅
5   T2.insert(A[1])
6   for k := 2 to n do
7     if A[k] < A[k - 1] then
8       T1.insert(A[k])
9       A[k] := T2.max()
10    else
11      T2.insert(A[k])
12    endif
13  endwhile

```

- (a) (4 points) Analyse the time consumption of the algorithm.
- (b) (4 points) Argue, using an invariant, that after the program is run, the resulting array  $A$  is sorted.

teht.	1	2	3	4	5	yht.
max.	8	10	10	8	8	44
op.						